

An Efficient Dynamic Mechanism

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1 Purpose of the Paper

- Dynamic Setting with Private Info
- Construct an Efficient, BB, BIC mechanism
- Extension of AGV/Arrow (1979)
- Idea: charge an agent for the *change* in the opponent utilities induced by his report
- Conditions of Self-Enforcing Mechanism (in a Markov Setting)

2 A Motivating Example

Buyer, Seller trade $x_t \in [0, 1]$ each period $t = 1, 2$

Seller's cost is $\frac{1}{2\theta_S}(x_t)^2$, parameter $\theta_S \in [0, 1]$ is private info

Buyer's valuation is 1 in $t = 1$, and θ_B in $t = 2$

NB θ_B is learnt between periods 1 and 2

2.1 First best:

$$\max_{x_1, x_2} \left(x_1 - \frac{1}{2\theta_S}(x_1)^2 + \theta_B x_2 - \frac{1}{2\theta_S}(x_2)^2 \right)$$

$$\chi_1(\theta_S, \theta_B) = \theta_S; \quad \chi_2(\theta_S, \theta_B) = \theta_S \theta_B$$

2.2 Expected externality transfers to induce truthful revelation:

$$\gamma_i(\theta_i) = E_{\tilde{\theta}_{-i}} U_{-i}(\theta_i, \tilde{\theta}_{-i})$$

$$\gamma_S(\theta_S) = E_{\tilde{\theta}_B} \left[\chi_1(\theta_S) + \tilde{\theta}_B \chi_2(\theta_S, \tilde{\theta}_B) \right] = \theta_S \left(1 + E_{\tilde{\theta}_B} \left[(\tilde{\theta}_B)^2 \right] \right)$$

$$\gamma_B(\theta_B) = E_{\tilde{\theta}_S} \left[-\frac{1}{2\tilde{\theta}_S} \left(\chi_1(\tilde{\theta}_S) \right)^2 - \frac{1}{2\tilde{\theta}_S} \left(\chi_2(\tilde{\theta}_S, \theta_B) \right)^2 \right] = -\frac{1}{2}(1 + (\theta_B)^2) E_{\tilde{\theta}_S} \left[\tilde{\theta}_S \right]$$

To balance the budget, take $\psi_i(\theta_i, \theta_{-i}) = \gamma_i(\theta_i) - \gamma_{-i}(\theta_{-i})$

2.3 The Problem:

B cannot make his report before $t = 1$

By the moment he learns θ_B , he has already observed the reported θ_S from trade, $\chi_1(\theta_S) = \theta_S$

B's utility in the mechanism:

$$\begin{aligned} \underset{\hat{\theta}_B}{\operatorname{argmax}} U_B(\theta_B, \hat{\theta}_B) &= \underset{\hat{\theta}_B}{\operatorname{argmax}} \left[\theta_B \chi_2(\theta_S, \hat{\theta}_B) + \gamma_B(\theta_B) \right] \\ &= \underset{\hat{\theta}_B}{\operatorname{argmax}} \left[\theta_B \theta_S \hat{\theta}_B - \frac{1}{2}(1 + (\theta_B)^2) E_{\tilde{\theta}_S} \left[\tilde{\theta}_S \right] \right] = \frac{E[\tilde{\theta}_S]}{\theta_S} \theta_B \end{aligned}$$

Knowing that B's strategy, S will also prefer to lie (a.s.).

2.4 Solution:

Let B's incentive transfer depend on both reports:

$$\text{instead of } \psi_B(\theta_S, \theta_B) = E_{\tilde{\theta}_S} U_S(\tilde{\theta}_S, \theta_B) - E_{\tilde{\theta}_B} U_B(\theta_S, \tilde{\theta}_B)$$

let $\psi_B(\theta_S, \theta_B) = U_S(\theta_S, \theta_B) + h(\theta_S)$ – will give him correct incentives

$$\text{Seller's utility: } U_S(\theta_S, \theta_B) + \psi_S(\theta_S, \theta_B) = E_{\tilde{\theta}_B} \left[U_S(\theta_S, \tilde{\theta}_B) + U_B(\theta_S, \tilde{\theta}_B) \right]$$

$$\text{BB: } \psi_B(\theta_S, \theta_B) = U_S(\theta_S, \theta_B) - E_{\tilde{\theta}_B} \left[U_S(\theta_S, \tilde{\theta}_B) + U_B(\theta_S, \tilde{\theta}_B) \right]$$

$$\psi_S(\theta_S, \theta_B) = -U_S(\theta_S, \theta_B) + E_{\tilde{\theta}_B} \left[U_S(\theta_S, \tilde{\theta}_B) + U_B(\theta_S, \tilde{\theta}_B) \right]$$

The idea is to charge an agent for the *change* in the opponent utilities induced by his report.

3 Set-up

I agents

$t = 1, 2, \dots$ countable/finite number of periods

Each period t each agent i privately observes $\theta_{i,t} \in \Theta_{i,t}$

individual type at t	$\Theta_{i,t}$	history of individual types up to t	$\Theta_i^t = \prod_{\tau=1}^t \Theta_{i,\tau}$
types at t	$\Theta_t = \prod_{i=1}^I \Theta_{i,t}$	history of types up to t	$\Theta^t = \prod_{i=1}^I \Theta_i^t$
the whole type space	$\Theta = \prod_{t=1}^{\infty} \Theta_t$		$\Theta = \Theta^{\infty}$

In each period t each agent i takes private decision $x_{i,t} \in X_{i,t}$ and there is a public decision $x_{0,t} \in X_{0,t}$.

Denote, equivalently, $X_t = \prod_{i=0}^I X_{i,t}$, $X^t = \prod_{\tau=1}^t X_{\tau}$, $X = X^{\infty}$.

Distribution of types over Θ_t is given by $\nu_t | x^{t-1}, \theta^{t-1}$ – conditioning on all past decisions and types.

A decision plan is a measurable function $\chi : \Theta \rightarrow X$,

where each $\chi_t(\theta)$ represents the decision made at t .

Any decision plan χ uniquely determines probability distribution $\mu[\chi]$ over the type space Θ .

Definition 1 We have independent types if for each t , x^{t-1} , and θ^{t-1} , the probability measure $\nu_t|x^{t-1}, \theta^{t-1}$ over Θ_t can be written in the form: $\nu_t|x^{t-1}, \theta^{t-1} = \prod_{i=1}^I \nu_{i,t}|x_o^{t-1}, x_i^{t-1}, \theta_i^{t-1}$ for some probability measures $\nu_{i,t}$ over $\Theta_{i,t}$ that depend only on public decision history x_o^{t-1} and on agent i 's private history $(x_i^{t-1}, \theta_i^{t-1})$.

NB $\mu[\chi]$ need not be independent across agents (influence through public decisions).

Agent i 's Payoff:

$$\sum_{t=1}^{\infty} \delta^t (u_{i,t}(x^t, \theta^t) + y_{i,t})$$

denote $U_i(x, \theta) = \sum_{t=1}^{\infty} \delta^t u_{i,t}(x^t, \theta^t)$.

Definition 2 We have private values if each agent i 's utility $U_i(x, \theta)$ depends only on the public decisions x_0 and the agent's private history (x_i, θ_i) .

NB The agent's expected utility conditional on time- t history (x^t, θ^t) is allowed to depend on other agents' information.

4 Mechanism

A mechanism is described by 2 observationally measurable functions:

decision plan $\chi : \Theta \rightarrow X$,

uniformly bounded transfer plan $\psi : \Theta \rightarrow (R^I)^\infty$

Public decisions $\chi_0(\theta)$ are implemented directly,

prescribed private decisions $\chi_i(\theta)$ for agents $i \geq 1$ can be disobeyed.

Transfer $\psi_{i,t}(\theta)$ is made to agent i in period t and depends only on the current history θ^t (observational measurability), the total discounted payments in the mechanism can be calculated as:

$$\Psi_i(\theta) = \sum_{t=1}^{\infty} \delta^t \psi_{i,t}(\theta^t)$$

Definition 3 A mechanism is budget balanced, if $\sum_i \psi_{i,t}(\theta^t) \equiv 0$ for all t .

5 The Game

Each agent privately observes his private type $\theta_{i,t}$

Each agent makes a public report $\hat{\theta}_{i,t}$

Each agent takes a private decision $x_{i,t}$. The mechanism implements the public decision $x_{0,t} = \chi(\hat{\theta}^t)$ and the payments $y_{i,t} = \psi_{i,t}(\hat{\theta}^t)$ to each agent.

6 Strategies

Agent i 's reporting strategy $\beta_i : \Theta \times \Theta_i \times X_i \rightarrow \Theta_i$, where the report at time t is observationally measurable, in that $\beta_{i,t}(\hat{\theta}, \theta_i, x_i)$ depends only on the history of reports $\hat{\theta}^{t-1}$ and the agent's private history θ_i^t, x_i^{t-1} .

Agent i 's decision strategy $\alpha_i : \Theta \times \Theta_i \times X_i \rightarrow X_i$, where the decision at time t is, again, observationally measurable.

A given strategy (β_i, α_i) induces a strategic plan $(\bar{\beta}_i, \bar{\alpha}_i)$, where $\bar{\beta}_i(\theta)$ denotes the agent i 's reports and $\bar{\alpha}_i(\theta)$ his decisions from following the strategy given that his type is θ_i and his opponents report θ_{-i} .

The construction is recursive, $t = 1, 2, \dots$:

$$\begin{aligned}\bar{\beta}_{i,t}(\theta) &= \beta_{i,t} \left(\left(\bar{\beta}_i^{t-1}(\theta_i^{t-1}), \theta_{-i}^{t-1} \right), \theta_i^t, \bar{\alpha}_i^{t-1}(\theta_i^{t-1}) \right) \\ \bar{\alpha}_{i,t}(\theta) &= \alpha_{i,t} \left(\left(\bar{\beta}_i^t(\theta_i^t), \theta_{-i}^t \right), \theta_i^t, \bar{\alpha}_i^{t-1}(\theta_i^{t-1}) \right)\end{aligned}$$

Observational measurability implies that $\bar{\beta}_{i,t}(\theta)$ depends only on $(\theta_i^t, \theta_{-i}^{t-1})$ and $\bar{\alpha}_{i,t}(\theta)$ only on θ^t .

Definition 4 For a given decision plan χ , agent i 's strategy (β_i, α_i) is a truthful-obedient strategy, if the corresponding strategic plan $(\bar{\beta}_i, \bar{\alpha}_i)$ has $\bar{\beta}_i(\theta) = \theta_i$ and $\bar{\alpha}_i(\theta) = \chi_i(\theta)$ for all $\theta \in \Theta$.

NB Does not specify anything about reporting after a lie or disobedience. "If has been truthful and obedient so far, he'll continue to be truthful and obedient".

7 The Team Mechanism

Let χ^* be the efficient decision plan (assumed to exist) - it maximizes $E_{\tilde{\theta}}^{\mu[\chi]} \left[\sum_i U_i \left(\chi \left(\tilde{\theta} \right), \tilde{\theta} \right) \right]$.

The Team Mechanism consists of the efficient decision plan χ^* and transfers:

$$\psi_{i,t}^M(\theta) = \sum_{j \neq i} u_{j,t}(\chi^*(\theta), \theta).$$

The resulting total discounted transfers are:

$$\Psi_i^M(\theta) = \sum_{j \neq i} U_j(\chi^*(\theta), \theta)$$

Proposition 1 Assume private values. Then for any mechanism with an efficient decision plan χ^* and total discounted transfers $\Psi_i^M(\theta)$ in the Team Mechanism, any truthful-obedient strategy profile is a BNE.

Proof. Define a new decision plan and transfers to incorporate a deviation to (β_i, α_i) :

$$\begin{aligned}\widehat{\chi}(\theta) &= (\chi_{-i}(\bar{\beta}_i(\theta), \theta_{-i}), \bar{\alpha}_i(\theta)) \\ \widehat{\psi}(\theta) &= \psi(\bar{\beta}_i(\theta), \theta_{-i}) \quad \text{and} \quad \widehat{\Psi}_i(\theta) = \sum_{t=1}^{\infty} \delta^t \widehat{\psi}_{i,t}(\theta)\end{aligned}$$

A truthful-obedient strategy in the new mechanism $(\widehat{\chi}, \widehat{\psi})$ will induce the same decisions and transfers as deviation (β_i, α_i) in the original mechanism (χ, ψ) . The respective payoff is $E_{\tilde{\theta}}^{\mu[\widehat{\chi}]} \left[U_i(\widehat{\chi}(\tilde{\theta}), \tilde{\theta}) + \widehat{\Psi}_i(\tilde{\theta}) \right]$.

For the Team Mechanism the new transfers are

$$\begin{aligned}\widehat{\Psi}_i(\theta) &= \sum_{j \neq i} U_j(\chi_{-i}^*(\bar{\beta}_i(\theta), \theta_{-i}), \chi_i^*(\bar{\beta}_i(\theta), \theta_{-i}), \bar{\beta}_i(\theta), \theta_{-i}) \\ &= \sum_{j \neq i} U_j(\widehat{\chi}_{-i}(\theta), \chi_i^*(\bar{\beta}_i(\theta), \theta_{-i}), \bar{\beta}_i(\theta), \theta_{-i}) = \sum_{j \neq i} U_j(\widehat{\chi}(\theta), \theta)\end{aligned}$$

Thus agent i's expected payoff from deviation is:

$$E_{\tilde{\theta}}^{\mu[\widehat{\chi}]} \left[U_i(\widehat{\chi}(\tilde{\theta}), \tilde{\theta}) + \sum_{j \neq i} U_j(\widehat{\chi}(\tilde{\theta}), \tilde{\theta}) \right].$$

Recall that χ^* maximizes this expression, and $\widehat{\chi} = \chi^*$ if (β_i, α_i) is a truthful-obedient strategy. Thus there is no profitable deviation from truthful-obedience. ■

8 Balancing the Transfers

So far, transfers of the Team Mechanism are not balanced.

Recall $\psi_{i,t}^M(\theta) = \sum_{j \neq i} u_{j,t}(\chi^*(\theta), \theta)$ $\Psi_i^M(\theta) = \sum_{j \neq i} U_j(\chi^*(\theta), \theta)$.

Construct balanced transfers for a mechanism (χ, ψ) :

$$\psi_{i,t}^B(\theta^t) = \gamma_{i,t}(\theta_i^t, \theta_{-i}^{t-1}) - \frac{1}{I-1} \sum_{j \neq i} \gamma_{j,t}(\theta_j^t, \theta_{-j}^{t-1}), \text{ where}$$

$$\gamma_{i,t}(\theta_i^t, \theta_{-i}^{t-1}) = \delta^{-t} \left(E_{\tilde{\theta}}^{\mu[\chi]|\theta_i^t, \theta_{-i}^{t-1}} \left[\Psi_i(\tilde{\theta}) \right] - E_{\tilde{\theta}}^{\mu[\chi]|\theta^{t-1}} \left[\Psi_i(\tilde{\theta}) \right] \right)$$

Proposition 2 Consider any mechanism (χ, ψ) in which there is a truthful-obedient BNE. If the agents' types are independent, then there is also a truthful-obedient BNE in the balanced mechanism (χ, ψ^B) .

Corollary 1 With independent types and private values, there is a truthtelling BNE in the Balanced Team Mechanism.

9 Markov games with Private States

A Folk-theorem like result for Markov games: efficiency can be sustained in equilibrium when the discount factor is close enough to 1.

9.1 The Game

Each agent i privately observes his private type $\theta_{i,t}$

Each agent i makes a public report $\hat{\theta}_{i,t}$

Each agent i chooses a publicly observable action $x_{0,i,t} \in X_{0,i,t}$, a privately observed action $x_{i,t} \in A_{i,t}$, and a publicly observable payment $z_{i,j,t} \geq 0$ to each agent j .

Let $x_t = ((x_{0,1,t}, \dots, x_{0,I,t}), x_{1,t}, \dots, x_{I,t})$,

$$y_{i,t} = \sum_{j=1}^I (z_{j,i,t} - z_{i,j,t})$$

then the payoff of agent i is, as before: $\sum_{t=1}^{\infty} \delta^t (u_{i,t}(x^t, \theta^t) + y_{i,t})$

9.2 Assumptions

•(a) $\Theta_t = \bar{\Theta}$ and $X_t = \bar{X}$ for all t , and all sets are finite.

(b) Per-period utilities $u_{i,t}(x^t, \theta^t)$ depend only on period t variables and do not depend on t , so $u_i(x_t, \theta_t)$.

- (c) Similarly, the probability measures $\nu_t|x^{t-1}, \theta^{t-1} = \nu|x_{t-1}, \theta_{t-1}$.
- (d) Independent types, private values.
- (e) $\exists \hat{x}_t \in X_t$ such that $u_i(x_{i,t}, \hat{x}_{-i,t}, \theta_t) = 0$ for all i , all $x_{i,t} \in \bar{X}$, and all $\theta_{i,t} \in \bar{\Theta}$.

Definition 5 *Decision policy χ is a Markov policy if each $\chi_t(\theta^t)$ depends only on the current state and does not depend on t directly, and thus can be written as $\chi(\theta_t)$. A decision policy χ^* is a Blackwell policy, if it is a Markov policy and there exists $\bar{\delta} < 1$ such that for all $\delta \in [\bar{\delta}, 1]$ it achieves a weakly greater present expected total surplus starting at any state than any other observationally measurable policy.*

Puterman (1994) proves existence and constructs a Blackwell policy.

In the Markov game the Team transfers can be written as:

$$\psi_{i,t}(\theta_t) = \sum_{j \neq i} u_j(\chi(\theta_t), \theta_t)$$

and the Balanced transfers can be written as:

$$\psi_{i,t}^B(\theta_{i,t}, \theta_{t-1}) = \gamma_i(\theta_{i,t}, \theta_{t-1}) - \frac{1}{I-1} \sum_{j \neq i} \gamma_j(\theta_{j,t}, \theta_{t-1})$$

$$\gamma_i^B(\theta_{i,t}, \theta_{t-1}) = \sum_{\tau=t}^{\infty} \delta^{\tau-t} \sum_{j \neq i} \left[E_{\tilde{\theta}}^{\mu_t^i[\chi]} | \theta_{i,\tau}, \theta_{\tau-1} u_j(\chi(\tilde{\theta}_\tau), \tilde{\theta}_\tau) - E_{\tilde{\theta}}^{\mu_t[\chi]} | \theta_{\tau-1} u_j(\chi(\tilde{\theta}_\tau), \tilde{\theta}_\tau) \right]$$

Definition 6 Given a Markov process, a set of states $S \subseteq \bar{\Theta}$ is called ergodic, if, starting in S , all system remains there with probability 1, and this is not true for any proper subset of S .

Proposition 3 Consider dynamic game satisfying assumptions 1-5. Let χ^* be a Blackwell decision policy for the game, and suppose that the Markov process induced by χ^* has a single ergodic set, and that its invariant distribution yields a positive expected total surplus. Then there exists $\delta^* < 1$ such that, for all $\delta \in (\delta^*, 1)$, the game has an efficient PBE that sustains decision plan χ^* .

10 Extensions

- Relax Markovian properties, look at limited dependence
- Allocations other than efficient ones
- Costly computation and communication
- Correlated values (efficiency with BB may be achievable more easily)

11 Summary

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- Conditions of Self-Enforcing Mechanism (in a Markov Setting)